

## Random Processes : Solutions 2

$$2.1 \text{ prob. of a "1"} = p = \frac{1}{6}$$

$$\text{prob. of "not 1"} = q = \frac{5}{6}.$$

prob. of "1" on k'th roll (following

$$(k-1 \text{ failures}) = q \times q \times q \dots \times q \times p = q^{k-1} p \quad (A)$$

i.e., the geometric distribution.

Expected # attempts:

$$\langle k \rangle = \sum_{k=1}^{\infty} k P_k = \sum_{k=1}^{\infty} k q^{k-1} p$$

$$= p \frac{d}{dq} \sum_{k=1}^{\infty} q^k = p \frac{d}{dq} \left( \frac{q}{1-q} \right)$$

$$= p \left( \frac{1}{1-q} + \frac{q}{(1-q)^2} \right) = \frac{p}{(1-q)^2} = \frac{1}{p} \quad (B)$$

In this case,  $\langle k \rangle = 6$ .

2.2 Derivation of required formulas is the same as above: results (A) and (B).

2.3 (a) The number of ways of drawing 6 balls from 49 is

$$\frac{49!}{6!(49-6)!} = \binom{49}{6}.$$

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(b) The machine can select r of "my"  $\checkmark$ 's in  $\binom{6}{r}$  ways, &  $6-r$  of the other 43 #'s in  $\binom{43}{6-r}$  ways.

Choices are independent (they are from different sets!), so total

$$\# \text{ ways to get } r \text{ right and } (6-r) \text{ wrong} = \underline{\underline{\binom{6}{r} \times \binom{43}{6-r}}}.$$

(c) Prob. of getting exactly r right is just the ratio (b)/(c), assuming that all possibilities are equally likely:

$$Pr = \binom{6}{r} \binom{43}{6-r} / \binom{49}{6}$$

(d) [Q] It's like 1:4 with  $M = 6$ ,  $N = 43$  and  $n = 6$ :

$$\binom{49}{6} = \binom{6}{0} \binom{43}{6} + \binom{6}{1} \binom{43}{5} + \dots + \binom{6}{6} \binom{43}{0}$$

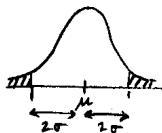
Dividing each side by  $\binom{49}{6}$  and using result (c) gives  $1 = \sum_{r=0}^6 Pr_r$ .

r	Pr	$\langle N_r \rangle = N P_r$
3	0.0177	577.883
4	$9.69 \times 10^{-4}$	317.13
5	$1.84 \times 10^{-5}$	6.04
6	$7.15 \times 10^{-8}$	2.3

So, not bad agreement?

(f) [We'll revisit this in Examples 4.]

$$2.4 \quad 25 \text{ cm} = \mu + 2\sigma$$



We know that  
prob of a result being  
more than  $2\sigma$  from mean is 5% for a  
Gaussian random variable. From the symmetry  
about  $\mu$ ,  $P(\text{result} > \mu + 2\sigma) = 2.5\%$ .

2.5 This is re-hashing  $2.1/2.2$ . In time  $t$ ,  
the aptd. makes  $k = t/\Delta t$  escape  
attempts. With  $p = \lambda \Delta t$ ,

$$\langle k \rangle = \langle t/\Delta t \rangle = \frac{1}{p} = \frac{1}{\lambda \Delta t},$$

$$\Rightarrow \langle t \rangle = \underline{\underline{\frac{1}{\lambda}}}.$$

Probability of survival to time  $t$

$$= q^{k-1} = (1-p)^{k-1}$$

$$= (1 - \lambda \Delta t)^{k-1}$$

$$\approx (1 - \frac{\lambda t}{k})^k \quad \begin{array}{l} \text{(we can ignore} \\ \text{this in comparison to} \\ k, \text{ since } (1 - \lambda \Delta t) \rightarrow 1 \\ \text{as } \Delta t \rightarrow 0. \end{array}$$

with  $k \rightarrow \infty$  as  $\Delta t \rightarrow 0$ . Use  $\ln(1-x) = -x - \frac{x^2}{2} - \dots$

$$\text{for } x = \frac{\lambda t}{k} \ll 1 :$$

$$\text{prob.} = e^{k \ln(1 - \lambda t/k)} = e^{k(-\frac{\lambda t}{k} - \frac{\lambda^2 t^2}{2k^2} - \dots)}$$

$$= e^{-\lambda t - \frac{\lambda^2 t^2}{2k} - \dots} \rightarrow \underline{\underline{e^{-\lambda t}}} \text{ for } k \rightarrow \infty$$

leads to 0 for  $k \rightarrow \infty$ .

2.6 Suppose that we have a total of  $N$  atoms today; then for each isotope,

$$N_{238} = 0.993 N$$

$$N_{235} = 0.007 N$$

If the numbers 1.8 Gy ago were  $N_{238}^0$  and  $N_{235}^0$ , then

$$N_{238}^0 = 2^{1.8/4.5} N_{238} = 2^{1.8/4.5} \times 0.993 N$$

$$N_{235}^0 = 2^{1.8/0.7} N_{235} = 2^{1.8/0.7} \times 0.007 N$$

[i.e., greater by factors  $2^{t/t_{1/2}}$  in each case.]

$\therefore$  Proportion of  $^{238}\text{U}$  1.8 Gy ago

$$\text{was } \frac{N_{238}^0}{N_{238}^0 + N_{235}^0} = \frac{2^{1.8/4.5} \times 0.993}{2^{1.8/4.5} \times 0.993 + 2^{1.8/0.7} \times 0.007}$$

$$= 0.969$$

$$\Rightarrow 96.9\% \text{ } ^{238}\text{U} \text{ and } 3.1\% \text{ } ^{235}\text{U}.$$